

# DIOPHANTINE EQUATIONS OF THE PILLAI TYPE: EXTENSIONS & APPLICATIONS

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Catalan problem:  $a^x - b^y = 1$

$3^2 - 2^3 = 1$  is the only solution with  $a, b, x, y \in \mathbb{N} \setminus \{1\}$   
[Mihalescu 2004]

Pillai problem:  $(P) \quad a^x - b^y = c \quad (c \in \mathbb{N})$

[widely open in general]

Partial results:

(i) abc-conjecture  $\Rightarrow (P)$  has only finitely many solutions

(ii)  $3^x - 2^y = c$  has at most 1 solution

[Stroeker-Tijdeman]

(iii)  $a^x - b^y = c$  ( $a, b, c$  fixed) has at most 2 solutions

[Bennett 2001]

(iv) Extensions to diophantine equations in finitely generated domains

[Evertse, Györy]

Koymans

# DIOPHANTINE EQUATIONS INVOLVING LINEAR RECURRENCES

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$$a^x \rightsquigarrow U_n \qquad b^y \rightsquigarrow V_m$$

$U_n, V_m$  L.r.s. of order  $k$

$$G_n = a_1 G_{n-1} + \dots + a_k G_{n-k}$$

$$x^k - a_1 x^{k-1} - \dots - a_k \quad \text{char. Pol. with integer coeff.}$$

dominating characteristic root  $\alpha$  :

$$|\alpha_1| > |\alpha_2| \geq \dots \geq |\alpha_k|$$

Problem PR  $U_n - V_m = c$

There exists a finite (and effectively computable) set  $\mathcal{C}$  such that PR has at least 2 solutions  $(n, m)$  if and only if  $c \in \mathcal{C}$ .

[Chim - Pink - Ziegler]

Survey Lecture summer 2022 :

discussion of many special cases

e.g.  $U_n =$  Fibonacci numbers, Tribonacci  
 $V_m = 2^m, 3^m$

# RECURRENCES AND POWER SEQUENCES

$U_n \dots$  binary linear r. s.

$$U_n - p_1^{x_1} \dots p_s^{x_s} = c \quad (p_i \dots \text{prime numbers})$$

has at most  $s$  solutions for  $c > c^+$   
(effect. comp.)

has at most  $s+1$  solutions for  $c < c^-$

[Ziegler, Debrecen 2022]

$$U_n = a_1 \alpha_1^n + a_2 \alpha_2^n + \dots + a_k \alpha_k^n \quad (\alpha > 1 \text{ irrational dominant root})$$

$a > 0$ ;  $a, \alpha$  mult. independent

and techn. cond.:  $\alpha^z - 1 = \alpha^x \alpha^y$  has no solutions  
 $z \in \mathbb{N}$ ;  $x, y \in \mathbb{Q}$ ;  $-1 < x < 0$

Then  $U_n - b^m = c$  has at most 2 solutions  $(n, m)$  for  $b > B, n > N_0$

(effect. computable)

[Heintze, Ti, Vukusic, Ziegler]

## Remarks

- ) techn. cond. can be algorithm. checked for any given recurrence  $(U_n)$
- ) without techn. cond. at most 3 solutions

## Several Examples and Problems

o)  $U_n$  Fibonacci sequence

$U_n$  Tribonacci sequence  $U_{n+3} = U_{n+2} + U_{n+1} + U_n$

$$U_{n+2} = U_{n+1} + 3U_n \quad (U_0 = 0, U_1 = 1)$$

o) In the Tribonacci case we found the following pairs  $(b, c)$  with 2 solutions:

$$(2, -8), (2, -3), (2, -1), (2, 0), (2, 5)$$

$$(3, -2), (3, 4), (5, -121), (5, -1), (5, 19)$$

$$(7, -5), (17, -15), (54, 220), (641, -137)$$

Are there further pairs  $(b, c)$  such that there are at least 2 solutions? Are there cases with at least 3 solutions?

## Tools for the proofs

- o) Linear forms in logarithms [Mateev, Laurent]
- o) Algebraic computations
- o) LLL-reduction

# PILLAI'S DENSITY PROBLEM

PILLAI'S CONJECTURE: For given  $c \in \mathbb{N}$ ,  
the equation  $(P) \quad a^n - b^m = c$  has at most 1  
solution  $(a, b, n, m)$  in integers  $\geq 2$ .

PILLAI'S THEOREM: For fixed  $a, b$  there  
is  $c_0 = c_0(a, b)$  such that  $\forall c > c_0$  the equation  $(P)$   
has at most 1 solution and furthermore

$$\begin{aligned} * \{c \in [1, x] : c = a^n - b^m \text{ for some } n, m \in \mathbb{N}\} &\sim \\ &\sim \frac{(\log x)^2}{2 \log a \cdot \log b} \quad (x \rightarrow \infty) \end{aligned}$$

LINEAR RECURRENCES  $U_n, V_m$  with dominating  
roots  $\alpha, \beta$  (resp.) replacing powers  $a^n, b^m$ .

Assuming  $\alpha, \beta$  multiplicatively independent:

$$\begin{aligned} * \{c \in [-x, x] : c = U_n - V_m \text{ for some } n, m \in \mathbb{N}\} &\sim \\ \sim * \{(n, m) \in \mathbb{N}^2 : |U_n - V_m| \leq x\} &\sim \\ \sim \frac{(\log x)^2}{\log |\alpha| \log |\beta|} &\quad (\text{as } x \rightarrow \infty) \end{aligned}$$

EQUIVALENT RESULT for multiplicatively independent algebraic numbers,  $|\alpha|, |\beta| > 1$

$$T_{\alpha, \beta}(x) \sim \frac{(\log x)^2}{\log |\alpha| \cdot \log |\beta|} \quad \left\{ (n, m) \in \mathbb{N}^2 : |\alpha^n - \beta^m| \leq x \right\}$$

QUESTION:  $\alpha, \beta$  transcendental?

COUNTER EXAMPLE:  $c = \sum_{i=0}^{\infty} 10^{-a(i)}$

extremely well approximable  
Liouville number

with  $a(0) = 1$   
 $a(i+1) = 10^{a(i)}$

$$\beta = 2, \alpha = 2^c \Rightarrow \frac{\log \alpha}{\log \beta} = c \notin \overline{\mathbb{Q}}$$

and  $|\alpha^n - \beta^m| \leq 1$  has infinitely many solutions  $(n, m) \in \mathbb{N}^2$ .

METRIC RESULT for Lebesgue almost all  $(\alpha, \beta)$  in  $\mathbb{C}^2$  with  $|\alpha|, |\beta| > 1$ :

$$T_{\alpha, \beta}(x) \sim \frac{(\log x)^2}{\log |\alpha| \cdot \log |\beta|} \quad (x \rightarrow \infty)$$

with remainder terms  $O(\log x (\log \log x)^2)$  (above)

$O(\log x (\log \log x))$  (below)

## SPECIAL CASES :

$$(1) \quad \alpha, \beta \in \overline{\mathbb{Q}}, \quad |\alpha| > 1, \quad \mu > 0 \text{ real}, \quad \beta = e^{\mu}$$

Then

$$T_{\alpha, \beta}(x) \sim \frac{(\log x)^2}{\mu \log |\alpha|} \quad (x \rightarrow \infty)$$

$$(2) \quad \mu_1, \mu_2 > 0 \text{ real algebraic and } \mathbb{Q}\text{-lin. indep.}, \\ \alpha = e^{\mu_1}, \quad \beta = e^{\mu_2}.$$

Then

$$T_{\alpha, \beta}(x) \sim \frac{(\log x)^2}{\mu_1 \cdot \mu_2} \quad (x \rightarrow \infty)$$

with remainder terms as above,  $\theta$ -constants are effective.

TOOL FOR PROOF: Waldschmidt's inhomogeneous linear form theorem

# WALDSCHMIDT'S THEOREM

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Let  $\mu_1, \dots, \mu_t \in \overline{\mathbb{Q}} \setminus \{0\}$  and let  $\log \mu_1, \dots, \log \mu_t$  be  $\mathbb{Q}$ -lin. indep and let  $\beta_0, \beta_1, \dots, \beta_t \in \overline{\mathbb{Q}}$ , not all = 0. Then

$$\begin{aligned} \log |\beta_0 + \beta_1 \log \mu_1 + \dots + \beta_t \log \mu_t| &\geq \\ &\geq -2^{t+25} t^{3t+9} (1 + \log k) \log B k^{t+2} \log A_1 \dots \log A_t, \end{aligned}$$

where  $k$  is degree of  $\mathbb{Q}(\mu_1, \dots, \mu_t, \beta_0, \dots, \beta_t)$  over  $\mathbb{Q}$

and

$$\log A_i \geq \max \left\{ h(\mu_i), \frac{e}{k} |\log \mu_i|, \frac{1}{k} \right\}$$

$$B \geq \max \left\{ k+1, \max_i k \log A_i, \max_i e^{h(\beta_i)} \right\}$$

logarithmic  
height

$$h(\eta) = \frac{1}{k} \left( \log |\alpha_0| + \sum_{i=1}^k \log (\max \{ |\eta^{(i)}|, 1 \}) \right)$$

$\eta^{(i)}$  conjugates of algebraic number  $\eta$



# PILLAI - TIJDEMAN EQUATION

$$Ax^m + By^m = Cx^n + Dy^n \quad \text{with } A, B \neq 0, |x| \neq |y|$$

$0 \leq n < m, Ax^m \neq Cx^n$

$$\Rightarrow m \in \bar{E} \quad (\text{effectively computable})$$

[ Shorey - Tijdeman ]

LUCAS SEQUENCES  $(U_n)_{n \geq 0}, U_0 = 0, U_1 = 1$

$$U_{n+2} = rU_{n+1} + U_n \quad (r \in \mathbb{N})$$

special case  $r=1$  Fibonacci sequence

$$AU_n + BU_m = CU_{n_1} + DU_{m_1} \quad \text{with } A, B \neq 0,$$

$n > m \geq 0, n_1 > m_1 \geq 0$   
 $AU_n \neq CU_{n_1}$

Then  $r < 14 \max\{|A|, |B|, |C|, |D|\}$

[ Ddamulira - Luca - Ti ]

Proof based on elementary algebraic computations and the following properties of Lucas numbers:

•)  $\gcd(U_n, U_m) = U_{\gcd(n, m)}$

•)  $\alpha^{n-2} \leq U_n \leq \alpha^{n-1}$

•)  $\alpha \geq \phi = \frac{\sqrt{5} + 1}{2}$

complete solution possible for small values of  $r$

# BINOMIAL POLYNOMIALS

$$\binom{x}{n} = \frac{x(x-1)\dots(x-n+1)}{n!} \in \mathbb{Q}[x] \quad (n \geq 1)$$

are integer-valued, i.e. elements of

$$\text{Int } \mathbb{Z} = \{f \in \mathbb{Q}[x] : f(\mathbb{Z}) \subseteq \mathbb{Z}\}.$$

$\binom{x}{n}$  is irreducible in the commutative ring  $\text{Int } \mathbb{Z}$ ,

i.e.  $a=bc$  implies that  $b$  or  $c$  is a unit. An

element  $a$  in a domain  $D$  is called

absolutely irreducible if  $a^m$  factors uniquely into irreducible elements ( $\forall m \in \mathbb{N}$ ).

(•)  $\binom{x}{p}$  is absolutely irreducible (abs-irr)  $\forall p$  prime number

[Frisch-Nakato via a graph-theoretic criterion]

(•)  $\binom{x}{n}$  is abs-irr  $\forall n \in \mathbb{N}$

[Rissner-Windisch via diophantine number theory]  
J. Algebra 2021

(•) Assume that for  $n > 10$  the  $k$  consecutive numbers  $n, (n-1), \dots, n-k+1$  are composite, then one of these numbers has a prime factor  $p > 2k$ .

# BASIC

# FACTS

(i)  $\binom{x}{1} = x$  is abs-irr

(ii) The property

$$\binom{x}{n}^m = f \cdot g \Rightarrow f = \pm \binom{x}{n}^k, g = \pm \binom{x}{n}^l$$

implies that  $\binom{x}{n}$  is abs-irr.

(iii) Let  $n, m \geq 2$ ;  $f, g \in \text{Int } \mathbb{Z}$  with  $\binom{x}{n}^m = f \cdot g$ .

Then for  $0 \leq i \leq n-1$  there exist  $k_i, l_i$  with  $k_i + l_i = m$  and

$$f = \pm \prod_{i=0}^{n-1} \left( \frac{x-i}{n-i} \right)^{k_i}; g = \pm \prod_{i=0}^{n-1} \left( \frac{x-i}{n-i} \right)^{l_i}$$

holds.

[Newton interpolation polynomials]

$$(iv) v_p(f(s)) = \sum_{j=0}^{n-1} (v_p(s-j) - v_p(n-j)) k_j \geq 0$$

$$v_p(g(s)) = \sum_{j=0}^{n-1} (v_p(s-j) - v_p(n-j)) l_j \geq 0$$

for all  $s \in \mathbb{Z}$  and  $p$  prime, where

$v_p(w)$  denotes the  $p$ -adic valuation of  $w \in \mathbb{Q}$ .

# VALUATION MATRIX

For  $n \in \mathbb{N}$  some notation :

$$\mathcal{P}_n = \{p: 0 < p \leq n, p \text{ prime}\}$$

For  $p \in \mathcal{P}_n$  let  $0 \leq r_{n,p} < p$  be the uniquely determined integer such that  $n \equiv r_{n,p} \pmod{p}$ .

$$R_{n,p} = \begin{cases} \{1, 2\} & \text{if } n = 2^s \text{ with } s > 1, p = 2 \\ \{1, 2, 3, 4\} & \text{if } n = 9 \text{ and } p = 3 \\ \{r: 1 \leq r \leq p - r_{n,p} - 1\} & \text{else} \end{cases}$$

## Valuation Matrix

$$A_n = \left( v_p(n+r-j) - v_p(n-j) \right)_{\substack{p \in \mathcal{P}_n, r \in R_{n,p} \text{ rows} \\ 0 \leq j \leq n-1 \text{ columns}}}$$

(•) row sums  $\sum_{j=0}^{n-1} (v_p(n+r-j) - v_p(n-j)) = 0$

$$\Rightarrow \text{rank } A_n < n.$$

## Key Proposition

$$\text{rank } A_n = n-1 \Rightarrow \binom{x}{n} \text{ is abs-irr.}$$

# ANALYSIS OF THE BLOCK STRUCTURE

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$p$ -block  $B_{n,p}$  is  $|\mathcal{R}_{n,p}| \times n$  matrix

$$B_{n,p} = \left( v_p(n+r-j) - v_p(n-j) \right)_{\substack{r \in \mathcal{R}_{n,p} \\ 0 \leq j \leq n-1}}$$

investigation of leftmost  $p$  and rightmost  $p-1$  columns

Proposition 1. Let  $n \in \mathbb{N}$ ,  $p \in \mathcal{P}_n$ ,  $B_{n,p} = (b_{r,j})$ . Then for  $0 \leq j \leq p-1$ ,  $1 \leq r \leq p - r_{n,p} - 1$ :

$$b_{r,j} = \begin{cases} -v_p(n - r_{n,p}) & \text{if } j = r_{n,p} \\ v_p(n - r_{n,p}) & \text{if } j = r + r_{n,p} \\ 0 & \text{else} \end{cases}$$

for  $n - (p-1) \leq j \leq n-1$ ,  $1 \leq r \leq p - r_{n,p} - 1$ :

$$b_{r,j} = \begin{cases} 1 & \text{if } j = n - p + r \\ 0 & \text{else} \end{cases}$$

Proposition 2. Let  $P = \max \mathcal{P}_n$ . Then

$$\text{rank } A_n \geq 2P - n - 1.$$

Corollary 1

If  $n = P$ ,  $\text{rank } A_n \geq P - 1$ , thus  $\text{rank } A_n = n - 1 \Rightarrow \binom{x}{n}$  is abs.-irr.

# NUMBER-THEORETIC TOOLS FOR COMPOSITE $n$

Theorem. Let  $n > 10$  and  $2 \leq k < n - P$ . Then there exists a prime  $p > 2k$  which divides one of the numbers  $n, n-1, \dots, n-k+1$ .

Proof by case study.

(I) large  $n \geq 4021520$ .

Fact 1 (Bertrand's postulate). For  $n \geq 3$  there exists a prime  $p$  with  $n/2 < p < n$ .

Fact 2 (Quantitative Chebyshev theorem). For  $n \geq 2010760$  there exists a prime  $p$  with  $n < p < (1 + \frac{1}{16597})n$ .

Fact 3 (Laishram-Shorey, 2005). Let  $k \geq 2$  and  $n > \max \{ k+13, \frac{279}{262}k \}$ . Then the product

$$n(n+1) \dots (n+k-1)$$

has a prime factor  $p > 2k$ .

(II) small  $n < 4021520$  and  $k \geq 5$ .

Fact 4 (Laishram-Shorey, 2006) Let  $k < n < k + 1.9 \cdot 10^{10}$ . Then there exist pairwise distinct primes  $p_0, p_1, \dots, p_{k-1}$  with  $p_i \mid n-i$  ( $\forall i=0, \dots, k-1$ ), provided that  $n, n-1, \dots, n-k+1$  are composite numbers.

③ Remaining cases  $k=2, 3, 4$  for small  $n$ .

Fact 5. By Mihăilescu's result on the Catalan equation the only consecutive positive integers which are non-prime prime powers are 8 and 9.

Fact 6. The only non-prime prime powers  $p^x$  and  $q^y$  less than  $10^{18}$  with  $p^x - q^y = 2$  are 25 and 27.

[Numerical result for Pillai equation]

This yields a proof of the number-theoretic theorem  $\Rightarrow \binom{x}{n}$  is abs-irr.